From: Erik Arvidsson arv@google.com

Subject: Re: FW: Your notes from the Google meeting

Date: October 6, 2009 at 10:17 AM

To: Istvan Sebestyen istvan@ecma-international.org

Cc: TC39 e-TC39@ecma-international.org



Sorry for not sending you the notes earlier.

Dear TC39 members, feel free to point out any errors and clear omissions since these notes are very rough at best.

```
Opt-in versioning?
<script type="application/ecmascript; version=5">
"use strict";
</script>
<script type="application/javascript; version=1.9">
let
yield
</script>
version=preharmony
      ph1
      ph2
ephemeron/name/cacthall
Mark: Page level meta tag?
Brendan: We talked about this before
Mike Samuel: Inline code to make sure that using a script by URL/path works:
use version harmony
This is attractive
Allen: Do we need to add namespacing
Brendan: If we only add a dozen items then we might get away without
Mark: Modules
Cormac: Do we need a way to allow setting a limitation on version per
page. The example is that newer versions might add catchalls which
changes the security model.
Mark: Another way to tackle this is to document the constraints under
which future versions and extensions may extend the language.
Brendan: Use lexical scope for example
Waldemar:
double a:
double b.
```



```
...
if (a < b) ... true
if (a < b) ... false
if (a < b) ... true
}
```

Brainstorm/discussion about host objects etc

Allen: As long as the ES5 spec is followed and no host objects or extensions are present reading a property is guaranteed to return the last value that was set.

8,9: Object model Execution model

Mark: make functions be the link to the environment. THe objects are native but its methods might be host objects

Brendan: The idea of taming host objects is something we should pursue.

```
Thursday 2009-09-24

Promises in E:

def \ r := a.foo(b, c) \text{ // sync}
def \ p := a<-foo(b, c) \text{ // async, eventually do}
def \ x := when(x, q)->\{
... \ x ... \ q ...
} catch(ex) { // optional}
... \ ex ...
}
def \ p := race([a, b, c, ...])
def \ p := timeBomb(millis, ex)
def \ p := race(a<-foo(b, c), timeBomb(3000, 'oops'))
def \ p := when(timeOut(3000)) -> \{
...
}
```

Brendan: ES next 2-3 years June GA 2012. Feature freeze in May 2011 (20 months). Definitional interpreter

Mark: Ephemerons require new kernel state

Allen: Weak refs as well

Mark: As soon as we introduce visible collection we need to express that

Waldemar/Allen/Mark: That can be done in prose

W: Grammar needs to be tightly integrated

```
airiiriai rioodo to bo tigirtiy iritogratod
A: We need a mapping at least
Fresh let or not in for (let i = 0; ...; i++)? Consensus to not get a new var.
Mark:
const a = [];
for (let i = 0; i < 3; i++) {
a.push(function() [{ return i * i; });
a[0](); // 4
Brendan:
for (let i in o) {
a.push(...)
for (const i in o) {
a.push(...)
(for (var i = E in o) \{...\} is valid in ES today)
No consensus after all?
Rob:
for (let x = []; x.length < 3; x.push(42)) {
Specify iteration order for ESH
Mark: Generator and finally?
Brendan: This has been solved in Python and Spidermonkey
Mark: return to label?
Brendan: Not without lambdas
Brendan: Maciej objected on the mailing list. Probably due to
implementation issues.
Allen: It is easy to implement
Brendan:
function gen() {
while (...) {
 try {
   yield x;
 } finally {
```

a - apu()

```
g - gon(),
g.next();
g.throw(e);

Let

let x;
x is undefined

redeclaration of let should be forbidden

for (var k in keys(o))
for (var v in values(o))
for each (var x in anlter)
```

Ephemerons

Allen: Adds overhead to the GC since the ephemerons have to be handled in a second pass